
Quantification

From last time:

Definition: *Proposition*

A declarative sentence that is either true (**T**) or false (**F**), but not both.

- Tucson summers never get above 100 degrees
- If the doorbell rings, then my dog will bark.
- You can take the flight if and only if you bought a ticket
- But **NOT** $x < 10$

Why do we want use variables?

- Propositional logic is not always expressive enough
- Consider: “All students like summer vacation”
- Should be able to conclude that “If Joe is a student, he likes summer vacation”.
- Similarly, “If Rachel is a student, she likes summer vacation” and so on.
- Propositional Logic does not support this!

Predicates

Definition: Predicate (a.k.a. Propositional Function)

A statement that includes at least one variable and will evaluate to either **true** or **false** when the variable(s) are assigned value(s).

- Example:

$$S(x) : (-10 < x) \wedge (x < 10)$$

$$E(a, b) : a \text{ **eats** } b$$

- These are not complete!

Predicates

Definition: Domain (a.k.a. Universe) of Discourse

The collection of values from which a variable's value is drawn.

- Example:

$$S(x) : (-10 < x) \wedge (x < 10), x \in \mathbb{Z}$$

$$E(a, b) : a \text{ eats } b, a \in \mathbf{People}, b \in \mathbf{Vegetables}$$

- In This Class: Domains may **NOT** hide operators
 - OK: **Vegetables**
 - Not-OK: **Raw Vegetables** ($\mathbf{Vegetable} \wedge \neg \mathbf{Cooked}$)

Predicates

$$S(x) : (-10 < x) \wedge (x < 10), x \in \mathbb{Z}$$

$$E(a, b) : a \text{ eats } b, a \in \mathbf{People}, b \in \mathbf{Vegetables}$$

- Can evaluate predicates at specific values (making them propositions):
 - What is $E(\text{Joe}, \text{Asparagus})$?
 - What is $S(0)$?

Combining Predicates with Logical Operators

- In $E(a, b) : a$ eats b , $a \in$ people, $b \in$ vegetables, change the domain of b to “raw vegetables”.
- $E(a, b) : a$ eats b , $a \in$ people, $b \in$ vegetables
- $C(b) : b$ is cooked, $b \in$ vegetables
- Combining the two:
 - $E(a, b) \wedge \neg C(b)$, $a \in$ people, $b \in$ vegetables

Quantification

- Idea: Establish truth of predicates over sets of values.
- Two common generalizations:
 - Universal Quantification ($\forall x P(x), x \in D$)
 - Considers all values from the domain of discourse
 - Existential Quantification ($\exists x S(x), x \in D$)
 - Considers one or more values from the domain of discourse

Note: Do not use the books non-standard $\exists!x$ notation
("uniqueness quantifier", Rosen 8/e p.46)

Quantifications in Propositional Logic

- Universal Quantification

- $\forall x P(x), x \in D \equiv P(d_0) \wedge P(d_1) \wedge P(d_2) \wedge \dots$

- Existential Quantification

- $\exists x S(x), x \in D \equiv S(d_0) \vee S(d_1) \vee S(d_2) \vee \dots$

Evaluating Quantified Predicates

- Universal Quantification:

Universal quantification of $P(x)$, $\forall x P(x)$, is the statement “ $P(x)$ holds for all objects x in the domain of discourse”

$\forall x P(x)$ is true only when $P(x)$ is true for **every** x in the domain, and false otherwise.

- Example 1:

- $Q(x) : x = x^2, x \in \mathbb{R}$

$$Q(-1) = \mathbf{False},$$

- $\forall x Q(x), x \in \{-1, 0, 1\}?$

$$-1 \neq (-1)^2$$

- A value x for which $P(x)$ is false is a counterexample $\forall x P(x)$

Evaluating Quantified Predicates

- Universal Quantification:

Universal quantification of $P(x)$, $\forall x P(x)$, is the statement “ $P(x)$ holds for all objects x in the domain of discourse”

$\forall x P(x)$ is true only when $P(x)$ is true for **every** x in the domain, and false otherwise.

- Example 2:

- $P(x, y) : x + y$ is even, $x, y \in \mathbb{Z}$

- $\forall x \forall y P(x, y), x, y \in \mathbb{Z}^{odd}$ **True! (Easy to prove)**

Evaluating Quantified Predicates

- Existential Quantification:

Existential quantification of $P(x)$, $\exists x P(x)$, is “There exists an element x in the domain of discourse such that $P(x)$ ”

$\exists x P(x)$ is true if **at least one** element x in the domain such that $P(x)$ is true

- Example 1:

- $Q(x) : x = x^2, x \in \mathbb{R}$

- $\exists x Q(x), x \in \{-1, 0, 1\}$?

$$Q(0) = \mathbf{True},$$
$$0 \neq 0^2$$

Evaluating Quantified Predicates

- Existential Quantification:

Existential quantification of $P(x)$, $\exists x P(x)$, is “There exists an element x in the domain of discourse such that $P(x)$ ”

$\exists x P(x)$ is true if **at least one** element x in the domain such that $P(x)$ is true

- Example 2:

- $P(x, y) : x + y$ is even, $x, y \in \mathbb{Z}$

- $\exists x \exists y P(x, y)$, $x, y \in \mathbb{Z}^{odd}$

True! Universal quantifier covers existential

Examples: Converting from English to Quantified Predicates

Example: Universal Quantification

- Consider this conversational English statement:

Every actor in The Office is great.

- How can we express that statement in logic notation?

— —WARNING — —

Several **INCORRECT** versions follow...

Only the last version is correct!

Example: Universal Quantification

- Consider this conversational English statement:

Every actor in The Office is great.

- How can we express that statement in logic notation?

Let $P(x)$: Actor x is great, $x \in \mathbf{People}$

$\forall x P(x), x \in \mathbf{People}$

Stilted English: For every person x , x is a great actor.

Con conversationally: All people are great actors.

PROBLEM: This is not quite the desired meaning

IDEA: Let's focus the domain!

Example: Universal Quantification

- Attempt #2: Every actor in The Office is great.

Let $P(x)$: Actor x is great, $x \in$ People who act in The Office

$\forall x P(x), x \in$ People who act in the Office

Stilted English: For every person x who acts in The Office, x is a great actor.

Con conversationally: All actors in the Office are great.

PROBLEM: The domain has a hidden predicate

IDEA: Let's create a new predicate.

Example: Universal Quantification

- Attempt #3: Every actor in The Office is great.

Let $P(x)$: Actor x is great, $x \in \mathbf{People}$

Let $Q(x)$: x was in The Office, $x \in \mathbf{People}$

$\forall x (P(x) \wedge Q(x)), x \in \mathbf{People}$

Stilted English: For every person x , x is in The Office and x is a great actor.

Con conversationally: All people are in the Office and are great.

PROBLEM: If true, implies that all people are in The Office!

IDEA: Try a different compound predicate

Example: Universal Quantification

- Attempt #4: Every actor in The Office is great.

Let $P(x)$: Actor x is great, $x \in \mathbf{People}$

Let $Q(x)$: x was in The Office, $x \in \mathbf{People}$

$\forall x (Q(x) \rightarrow P(x)), x \in \mathbf{People}$

Stilted English: For every person x , if x was in The Office, then x is a great actor.

Con conversationally: Every actor in the Office is great.

PROBLEM: Isn't $P(x)$, really two predicates in one?

IDEA: Break it apart

[Why not $P(x) \rightarrow Q(x)$? That says all great actors are in the office]

Example: Universal Quantification

- Attempt #5: Every actor in The Office is great.

Let $P(x)$: x is an actor, $x \in \mathbf{People}$

Let $A(x)$: x is a great actor, $x \in \mathbf{People}$

Let $Q(x)$: x was in The Office, $x \in \mathbf{People}$

$\forall x ((Q(x) \wedge P(x)) \rightarrow A(x)), x \in \mathbf{People}$

Stilted English: For every person x , if x is an actor and was in The Office, then x is a great actor.

Con conversationally: Every actor in the Office is great.

— — **SUCCESS!** — —

(This is the version to learn!)

Implicit Quantification

- The “all” can be implicit in the English statement
- Example:
 - Adding an odd # to itself produces an even #

$O(x) : x \text{ is odd}, x \in \mathbb{R}$

$E(x) : x \text{ is even}, x \in \mathbb{R}$

$\forall x (O(x) \rightarrow E(x + x)), x \in \mathbb{Z}$

$\forall x (O(x) \rightarrow \overline{O(x + x)}), x \in \mathbb{Z}$

Note the implicit \forall is implicit in the sentence

Example: Existential Quantification

- Consider this conversational English statement:

At least one breed of dog is cute.

- How can we express that statement in logic notation?

Let $C(x)$: x is cute, $x \in$ Dog Breeds

$\exists x C(x), x \in$ Dog Breeds

English: There is at least one dog breed x
such that x is cute.

Example: Existential Quantification

- Express this more specific statement in logic:

Some of the large fluffy dog breeds are cute.

Let $L(x)$: x is large, $x \in$ Dog Breeds

Let $F(x)$: x is fluffy, $x \in$ Dog Breeds

Let $C(x)$: x is cute, $x \in$ Dog Breeds

$\exists x (L(x) \wedge F(x) \wedge C(x)), x \in$ Dog Breeds

These alternatives don't work! Why?

$\exists x (L(x) \wedge F(x)) \rightarrow C(x), x \in$ Dog Breeds

$\exists x (L(x) \wedge C(x)) \rightarrow F(x), x \in$ Dog Breeds

Example: Existential Quantification

- Express this statement in logic:

Every last one of the large fluffy dog breeds are cute.

Let $L(x)$: x is large, $x \in$ Dog Breeds

Let $F(x)$: x is fluffy, $x \in$ Dog Breeds

Let $C(x)$: x is cute, $x \in$ Dog Breeds

$\forall x [(L(x) \wedge F(x)) \rightarrow C(x)], x \in$ Dog Breeds

If a dog is both large and fluffy, then it is cute.

(vs. ... $\wedge C(x)$): **All** dog breeds are large, fluffy and cute

Typically \forall pairs with \rightarrow and \exists goes with \wedge

Nested Quantifiers

- Sometimes we may need to use multiple quantifiers
 - We can't express "Everybody loves someone" using a single quantifier.
 - Suppose predicate **loves**(x, y): "Person x loves person y "

Nested Quantifiers

- **loves**(x, y): “Person x loves person y ”
- The four possible nestings:

- $\forall x \forall y$ **loves**(x, y)
- $\exists x \exists y$ **loves**(x, y)

Same quantifiers

- $\exists x \forall y$ **loves**(x, y)
- $\forall x \exists y$ **loves**(x, y)

Mixed quantifiers

Evaluating Nested Quantifiers of the Same Type

- Example: **loves**(x, y): “Person x loves person y ”
 - $\forall x \forall y$ **loves**(x, y)
 - “Everyone loves everyone”.
 - $\exists x \exists y$ **loves**(x, y)
 - “There is someone who loves someone else” (or possibly themselves!)

Evaluating Mixed Quantifiers

- Example: **loves**(x, y): “Person x loves person y ”
 - $\exists x \forall y$ **loves**(x, y)
 - “There is someone who loves everyone”
 - $\forall x \exists y$ **loves**(x, y)
 - “Everyone loves at least one person (possibly themselves)”

Evaluating Mixed Quantified

- Distinguishing $\exists x \forall y S(x, y)$ from $\forall i \exists k T(i, k)$
 - $\exists x \forall y S(x, y)$ - “There exists an x such that, for every y , $S(x, y)$ is true.”
 - Somewhere in x 's domain is an x that can be paired with any y 's domain to make $S(x, y)$ true.
 - $\forall i \exists k T(i, k)$ - “For any i there exists a k such that $T(i, k)$ is true.”
 - No matter which i is selected, we can find some k to pair with the i to make $T(i, k)$ true. (Note that the k may vary with the i)

Evaluating Mixed Quantified

- Example:
 - Given $P(x, y) : x - y = 0, x, y \in \mathbb{Z}$
 - Evaluate: $\exists x \forall y P(x, y)$

Evaluating Mixed Quantified

- Example:
 - Given $P(x, y) : x - y = 0, x, y \in \mathbb{Z}$
 - Evaluate: $\exists x \forall y P(x, y) = \mathbf{False}$
 - (No such magical integer x exists!)

x	y	$P(x, y)$
0	\vdots	\vdots
0	\vdots	\vdots
0	\vdots	\vdots
0	-1	F
0	0	T
0	1	F
0	\vdots	\vdots
0	\vdots	\vdots
0	\vdots	\vdots

Evaluating Mixed Quantified

- Example:

- Given $P(x, y) : x - y = 0, x, y \in \mathbb{Z}$

- Evaluate: $\exists x \forall y P(x, y) = \mathbf{False}$

- Evaluate: $\forall x \exists y P(x, y) = \mathbf{True}$

- (No matter the x , there's an integer y ($y = x$) that makes $P(x, y)$ true.)

x	y	$P(x, y)$
-3	-3	T
-2	-2	T
-1	-1	T
0	0	T
1	1	T
2	2	T
3	3	T

Negation of Quantified Expressions

- Remember De Morgan's Laws for Propositions? Well...

Definition: Generalized De Morgan's Laws

The Generalized De Morgan's Laws are the pair of equivalences:

$$\neg \forall x P(x) \equiv \exists \neg P(x)$$

$$\neg \exists x P(x) \equiv \forall \neg P(x)$$

Demonstration: $\neg \forall x P(x) \equiv \exists x \neg P(x)$

- **Reminder:**

$$\forall x S(x), x \in D \equiv S(d_1) \wedge S(d_2) \wedge S(d_3) \dots$$

$$\exists x S(x), x \in D \equiv S(d_1) \vee S(d_2) \vee S(d_3) \dots$$

Demonstration: $\neg \forall x P(x) \equiv \exists x \neg P(x)$

- Let $S(x) : x$ is cute, $x \in D$. Let $D = \{\text{all dog breeds}\}$

$$\forall x S(x), x \in D \equiv S(d_1) \wedge S(d_2) \wedge S(d_3) \dots$$

$$\begin{aligned} \neg \forall x S(x), x \in D &\equiv \neg(S(d_1) \wedge S(d_2) \wedge S(d_3) \dots) \\ &\equiv \neg S(d_1) \vee \neg S(d_2) \vee \neg S(d_3) \dots \\ &\equiv \exists x \neg S(x), x \in D \end{aligned}$$

$$\exists x S(x), x \in D \equiv S(d_1) \vee S(d_2) \vee S(d_3) \dots$$

Negation of Nested Quantifiers

- Apply De Morgan's Laws from left to right
- Example:

$$\begin{aligned} & \neg(\exists x \forall y (G(x) \vee \neg H(y))) \\ \equiv & \forall x \neg(\forall y (G(x) \vee \neg H(y))) \quad \text{[General DeMorgan]} \\ \equiv & \forall x \exists y \neg(G(x) \vee \neg H(y)) \quad \text{[General DeMorgan]} \\ \equiv & \forall x \exists y (\neg G(x) \wedge H(y)) \quad \text{[DeMorgan]} \end{aligned}$$

Expressing “Exactly one...” Statements

- Consider the conversational (& correct!) English statement

Only one citizen of Montana is a member of the U.S. House of Representatives

- And consider this awkward but useful rewording:

A member of the US House of Representatives exists in the set of citizens of Montana, and, if anyone in Montana is a member of the House, that person is the representative

Expressing “Exactly one...” Statements

- That rewording is useful because it can be directly expressed logically:

$R(x)$: x is a member of the US House of Representatives, $x \in \mathbf{People}$

$\exists x(R(x) \wedge \forall y[R(y) \rightarrow (y = x)]), x, y \in \mathbf{Citizens of Montana}$

This domain should be simplified, but using it makes the logic easier to read (for now)

Expressing “Exactly one...” Statements

- That rewording is useful because it can be directly expressed logically:

$R(x)$: x is a member of the US House of Representatives, $x \in \mathbf{People}$

$\exists x(R(x) \wedge \forall y[R(y) \rightarrow (y = x)]), x, y \in \mathbf{Citizens\ of\ Montana}$

- Interpretation: (At least one) \wedge (No more than one)
- \therefore Impossible for there to be two representatives!

Expression “Exactly two...” Statement

- Key observation:

Exactly 2 \equiv At least 2 \wedge At most 2

$$(n = 2) \equiv (n \geq 2) \wedge (n \leq 2)$$

- Awkward English:

At least two citizens of Montana are U.S. Senators,
and at most two citizens of Montana are U.S.
Senators.

- Better:

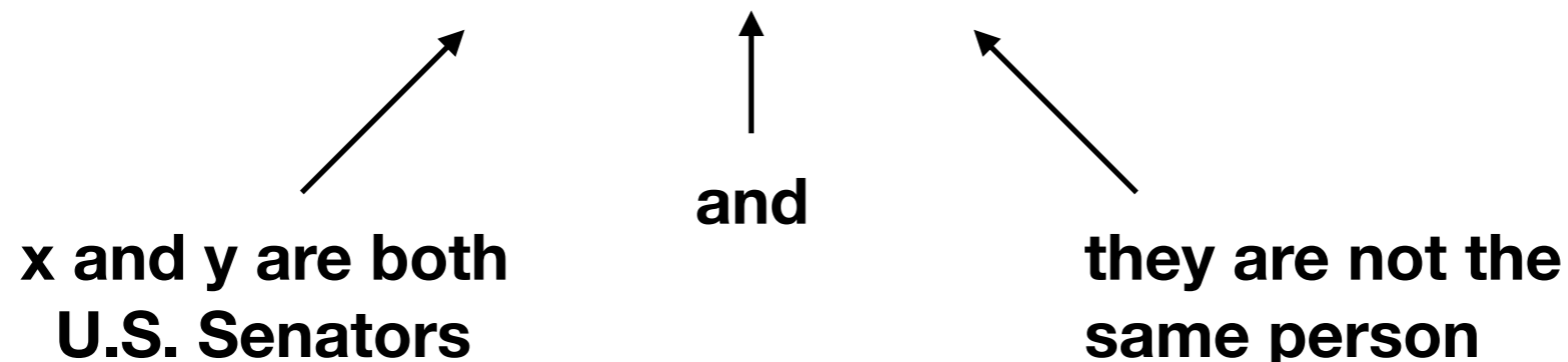
Exactly two citizens of Montana are U.S. Citizens

Expression “Exactly two...” Statement

- Consider the two halves separately:
 1. “At least two citizens of Montana are U.S. Senators”
 2. “At most two citizens of Montana are U.S. Senators”

Expressing “At least two”

- How do we express “At least two citizens of Montana are U.S. Senators”? Let $S(x) : x$ is a U.S. Senator, $x \in \text{People}$
- Why doesn't this work?
 - $\exists x \exists y (S(x) \wedge S(y)), x, y \in \text{Citizens of Montana}$ (**x, y could be identical**)
- Correct version:
 - $\exists x \exists y (S(x) \wedge S(y) \wedge (x \neq y)), x, y \in \text{Citizens of Montana}$



Expression “Exactly two...” Statement

- Consider the two halves separately:
 1. “At least two citizens of Montana are U.S. Senators”
 $S(x) : x \text{ is a U.S. Senator, } x \in \mathbf{People}$
 $\exists x \exists y (S(x) \wedge S(y) \wedge (x \neq y)),$
 $x, y \in \mathbf{Citizens of Montana}$
 2. “At most two citizens of Montana are U.S. Senators”

Expressing “At most two”

- How do we express “At most two citizens of Montana are U.S. Senators”? Let $S(x) : x$ is a U.S. Senator, $x \in \text{People}$

- Start with “at most one”:

- $\forall x \forall y (S(x) \wedge S(y)) \rightarrow (x = y)$

- Extended to 2 (i.e. “at most two”):

- $\forall x \forall y \forall z ((S(x) \wedge S(y) \wedge S(z)) \rightarrow$
 $((x = y) \vee (x = z) \vee (y = z)),$
 $x, y \in \text{Citizens of Montana}$

If x, y, z are all U.S. Senators...

... then there is **at least** one pair which are the same person

Expression “Exactly two...” Statement

- Consider the two halves separately:
 1. “At least two citizens of Montana are U.S. Senators”
 $S(x) : x \text{ is a U.S. Senator, } x \in \mathbf{People}$
 $\exists x \exists y (S(x) \wedge S(y) \wedge (x \neq y)),$
 $x, y \in \mathbf{Citizens of Montana}$
 2. “At most two citizens of Montana are U.S. Senators”
 $\forall x \forall y \forall z ((S(x) \wedge S(y) \wedge S(z)) \rightarrow$
 $(x = y \vee y = z \vee x = z))$
 $x, y, z \in \mathbf{Citizens of Montana}$

Expression “Exactly two...” Statement

- Finally, **AND** together

$$\exists x \exists y (S(x) \wedge S(y) \wedge (x \neq y))$$

- and

$$\forall x \forall y \forall z ((S(x) \wedge S(y) \wedge S(z)) \rightarrow (x = y \vee y = z \vee x = z))$$

Expression “Exactly two...” Statement

- Finally, **AND** together

$$\exists x \exists y (S(x) \wedge S(y) \wedge (x \neq y))$$

- and

$$\forall x \forall y \forall z ((S(x) \wedge S(y) \wedge S(z)) \rightarrow (x = y \vee y = z \vee x = z))$$

$$\begin{aligned} &\exists x \exists y (S(x) \wedge S(y) \wedge (x \neq y) \wedge \\ &\quad \forall z [S(z) \rightarrow (z = x \vee z = y)]), \\ &x, y, z \in \mathbf{Citizens\ of\ Montana} \end{aligned}$$

Why is the second half simplified?

Reminders

- Homework 2 due **this** Friday (06/19)
- Grades for Homework 1 should be posted before Friday
- Quiz 1 **this** Tuesday (on Logic)