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# Arguments

# Monty Python's "The Argument Clinic"

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## Definition: Argument

A connected series of statements to establish a definite proposition

### Featuring:

Michael Palin as "Man"  
Rita Davies as "Receptionist"  
Graham Chapman as "Mr. Barnard"  
John Cleese as "Mr. Vibrating"  
Eric Idle as "Complainer"  
Terry Jones as "Spreaders"



<https://www.youtube.com/watch?v=xpAvcGcEc0k>

# Inductive and Deductive Reasoning

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## **Definition:** Inductive Argument

An argument that moves from specific observations to general conclusions

## **Definition:** Deductive Argument

An argument that uses accepted general principles to explain a specific situation

# Inductive and Deductive Reasoning

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- Example:

- “Students who do well on the midterm do well in the class”

⇒ **Inductive**

- “Well-done hamburgers are safer to eat than medium-rare hamburgers”

⇒ **Deductive (heat kills bacteria)**

# Inductive and Deductive Reasoning

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- What type of argument is this?
  - 3 is a prime number, 5 is a prime number, and 7 is a prime number. Therefore all odd integers above 1 are prime numbers

⇒ **Inductive**

**(Specific examples lead to a general conclusion - that happens to be incorrect!)**

# Playposit Question

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Is the following argument inductive or deductive?

“People born before 2003 can vote”

- A. Inductive
- B. Deductive

# Structure of a Deductive Argument

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$$(p_1 \wedge p_2 \wedge \dots \wedge p_n) \rightarrow q$$

- (Recall the form: Hypothesis  $\rightarrow$  Conclusion)
- Approach: Logical principles are applied to the given hypothesis to see if the conclusion follows from them

- Common notations:

- $p_1 \wedge p_2 \wedge \dots \wedge p_n / \therefore q$  and

$$\frac{\begin{array}{c} p_1 \\ p_2 \\ \vdots \\ p_n \end{array} \text{Premises}}{\therefore q}$$

**We will use the second one**

# Valid and Sound Arguments

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## Definition: Valid Argument

Any deductive argument of the form

$(p_1 \wedge p_2 \wedge \dots \wedge p_n) \rightarrow q$  is valid if the conclusion must follow from the hypotheses

## Definition: Sound Argument

A valid argument that also has true premises



# Valid and Sound Arguments

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- Example:

**All men are mortal  
Socrates is mortal**

---

**∴ All men are Socrates**

**Not a valid argument!**

- Example:

**Socrates is a man  
All men are mortal**

---

**∴ Socrates is mortal**

**Valid! (Why?)**

# How can we show that arguments are valid?

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- Truth tables
  - What if our argument has 10 propositions in it?
- Rules of Inference
  - Building blocks to construct (or validate) complicated arguments
  - Always valid no matter the premises and conclusion (but only sound when premises are true!)

- Form: 
$$\begin{array}{c} p_1 \\ p_2 \\ \vdots \\ p_n \\ \hline \therefore q \end{array}$$

Corresponding propositional logic statement,  $(p_1 \wedge p_2 \wedge \dots \wedge p_n) \rightarrow q$ , is always a [tautology](#)!

# Rules of Inference

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## 1. Modus Ponens (Method of affirming):

$$\frac{p \quad p \rightarrow q}{\therefore q}$$

**English:** “If  $p$ , and  $p$  implies  $q$ , then  $q$ ”

**Propositional Logic:**  $(p \wedge (p \rightarrow q)) \rightarrow q$

# Example: Modus Ponens

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- Example: “When Rodger squeaks his toy, he wants to play. He is squeaking his toy.” Does Rodger want to play?

- $s$  : Rodger squeaks his toy,
- $p$  : Rodger wants to play

(1)  $s$  (Given)

(2)  $s \rightarrow p$  (Given)

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**Vertical structure leaves room for justification**

# Rules of Inference

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## 2. Modus Tollens (Method of Denying):

$$\frac{\neg q \quad p \rightarrow q}{\therefore \neg p}$$

**English:** “If it is not the case that  $q$ , and  $p$  implies  $q$ , then it must not be the case that  $p$ ”

**Propositional Logic:**  $(\neg q \wedge (p \rightarrow q)) \rightarrow \neg p$

# Example: Modus Tollens

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- Example: “When Rodger squeaks his toy, he wants to play. He does not want to play.” Is Rodger squeaking his toy?

- $s$  : Rodger squeaks his toy,
- $p$  : Rodger wants to play

- (1)  $\neg p$  (Given)
- (2)  $s \rightarrow p$  (Given)
- 



# Rules of Inference

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## 3. Hypothetical Syllogism (Transitivity of Implication):

$$\begin{array}{l} p \rightarrow q \\ q \rightarrow r \\ \hline \therefore p \rightarrow r \end{array}$$

**English:** “If  $p$  implies  $q$  and  $q$  implies  $r$ , then  $p$  must imply  $r$ .”

**Propositional Logic:**

$$((p \rightarrow q) \wedge (q \rightarrow r)) \rightarrow (p \rightarrow r)$$

# Example: Hypothetical Syllogism

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- Example: “If it’s spring, then there is pollen in the air. When there is pollen in the air, I sneeze.” If it’s spring, do I sneeze?
  - $s$  : It is spring
  - $p$  : There is pollen the air
  - $z$  : I sneeze

$$\begin{array}{lll} (1) & s \rightarrow p & (\text{Given}) \\ (2) & p \rightarrow z & (\text{Given}) \end{array}$$

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# Playposit Question

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Which rule is used in the following argument:

“I will cook if it is the weekend. I did not cook today.  
Therefore it is not the weekend”

- A. Modus Ponens
- B. Modus Tollens
- C. Hypothetical Syllogism

# Rules of Inference

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## 4. Disjunctive Syllogism (One or the Other):

$$\frac{p \vee q}{\neg p} \\ \hline \therefore q$$

**English:** “If  $p$  or  $q$  and it is not the case that  $p$ , then  $q$ .”

**Propositional Logic:**  $((p \vee q) \wedge \neg p) \rightarrow q$

## 5. Addition:

$$\frac{p}{\therefore p \vee q}$$

**English:** “If  $p$ , then  $p$  or  $q$ .”

**Propositional Logic:**  $p \rightarrow (p \vee q)$

# Rules of Inference

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## 6. Simplification:

$$\frac{p \wedge q}{\therefore p}$$

**English:** “If  $p$  and  $q$ , then  $p$ .”

**Propositional Logic:**  $(p \wedge q) \rightarrow p$

## 7. Conjunction:

$$\frac{p}{\therefore p \wedge q}$$
$$\frac{q}{\therefore p \wedge q}$$

**English:** “If  $p$ , and also  $q$ , then  $p$  and  $q$ .”

**Propositional Logic:**  $((p) \wedge (q)) \rightarrow (p \wedge q)$

# Rules of Inference

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## 8. Resolution:

$$\begin{array}{l} p \vee q \\ \neg p \vee r \\ \hline \therefore q \vee r \end{array}$$

**English:** “If  $p$  or  $q$ , and it is not the case that  $p$  or (it is the case that)  $r$ , then  $q$  or  $r$ ”

**Propositional Logic:**  $((p \vee q) \wedge (\neg p \vee r)) \rightarrow (q \vee r)$

- Why is this true?
  - If  $p$  is true, then  $\neg p$  is false. Thus, by the second premise,  $r$  must be true.
  - If  $p$  is false, then  $\neg p$  is true. Thus, by the first premise,  $q$  must be true
  - Thus, either  $q$  or  $r$  must be true.

# Example: Resolution

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- Example: “I walk my dog or it is raining. It is not raining or the wash is full of water ”.
  - $w$  : I walk my dog
  - $r$  : It is raining
  - $f$  : The wash is full

(1)             $w \vee r$             (Given)

(2)             $\neg r \vee f$             (Given)

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# Summary: Rules of Inference

## Learn These!

Name	Rule of Inference
Modus Ponens	$\frac{p \quad p \rightarrow q}{\therefore q}$
Modus Tollens	$\frac{p \rightarrow q \quad \neg q}{\therefore \neg p}$
Hypothetical Syllogism	$\frac{p \rightarrow q \quad q \rightarrow r}{\therefore p \rightarrow r}$
Disjunctive Syllogism	$\frac{p \vee q \quad \neg p}{\therefore q}$
Addition	$\frac{p}{\therefore p \vee q}$
Simplification	$\frac{p \wedge q}{\therefore p}$
Conjunction	$\frac{p \quad q}{\therefore p \wedge q}$
Resolution	$\frac{p \vee q \quad \neg p \vee r}{\therefore q \vee r}$

# Playposit

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Which rule of inference is used in the following argument:

“Norm eats chicken wings or lasagna. Norm doesn’t eat lasagna. Therefore, Norm eat’s chicken wings. “

A. Modus Ponens

E. Addition

B. Modus Tollens

F. Simplification

C. Hypothetical Syllogism

G. Conjunction

D. Disjunctive Syllogism

H. Resolution

# Using Rules of Inference

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- Example 1: Is the argument below valid?
  - If 191 is divisible by 7, then  $191^2$  is divisible by 49.
  - 191 is divisible by 7, so  $191^2$  must be divisible by 49.

**Yes! The supporting rule is Modus Ponens:**

*a* : 191 is divisible by 7

*b* :  $191^2$  is divisible by 49

(1)             $a \rightarrow b$     (Given)

(2)             $a$                     (Given)

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# Using Rules of Inference

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- Example 2: If my advisor sends me an email at 10pm, I have to work late. If he doesn't, then I will get plenty of sleep. If I get plenty of sleep, I'll be more productive tomorrow.
- Show that if I don't work late, I will be productive tomorrow.

**Begin by identifying propositions:**

**$p$  : My advisor sends me an email at 10pm**

**$q$  : I work late**

**$r$  : I get plenty of sleep**

**$s$  : I am productive tomorrow**

# Using Rules of Inference

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- Example 2: If my advisor  $p$  sends me an email at 10pm, I have to work  $q$  late. If he doesn't, then I will  $r$  get plenty of sleep. If I get plenty of sleep, I'll be more  $s$  productive tomorrow.
- Show that if I don't work late, I will be productive tomorrow.

**Next, identify the givens and the desired conclusion**

$$\begin{array}{l} \text{Givens} \\ \text{Conclusion} \end{array} \quad \begin{array}{l} p \rightarrow q \\ \neg p \rightarrow r \\ r \rightarrow s \\ \hline \therefore \neg q \rightarrow s \end{array}$$

**But how do we get there?**

# Using Rules of Inference

Propositions:	Givens and Conclusion:
$p$ : My advisor sends me an email at 10pm $q$ : I work late $r$ : I don't get much sleep $s$ : I am productive tomorrow	$\begin{array}{l} p \rightarrow q \\ \neg p \rightarrow r \\ r \rightarrow s \\ \hline \therefore \neg q \rightarrow s \end{array}$
(1) $p \rightarrow q$ (Given)	
(2) $\neg q \rightarrow \neg p$ (Contrapositive of (1))	
(3) $\neg p \rightarrow r$ (Given)	
(4) $\neg q \rightarrow r$ (Hypothetical Syllogism using (2) and (3))	
(5) $r \rightarrow s$ (Given)	
(6) $\therefore \neg q \rightarrow s$ (Hypothetical Syllogism using (4) and (5))	

**Note: This is a Formal Proof!**

# Using Rules of Inference

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- Example 3: We go hiking today or its over 90 degrees. It is under 90 degrees or we go to Andrew's to play games. We do not go hiking or we go to Andrew's to play games. We won't eat too many cookies only if we don't go to Andrew's to play games.
- Show that we eat too many cookies.

$p$  : We go hiking

$q$  : It is over 90 degrees

$r$  : We go to Andrew's house to play games

$s$  : We eat too many cookies

# Using Rules of Inference

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- **Example 3:** We go hiking  $p$  today or its over 90  $q$  degrees. It is under 90 degrees or we go to Andrew's  $r$  to play games. We do not go hiking or we go to Andrew's to play games. We won't eat too many cookies  $s$  only if we don't go to Andrew's to play games.
- Show that we eat too many cookies.

**Givens**

$$p \vee q$$

$$\neg q \vee r$$

$$\neg p \vee r$$

$$\neg s \rightarrow \neg r$$

**Conclusion**

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$$\therefore ?$$

# Using Rules of Inference

## Propositions:

$p$  : We go hiking

$q$  : It is over 90 degrees

$r$  : We go to Andrew's house to play games

$s$  : We eat too many of cookies

## Givens and Conclusion:

$$p \vee q$$
$$\neg q \vee r$$
$$\neg p \vee r$$
$$\neg s \rightarrow \neg r$$

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$$\therefore ?$$

- |     |                             |                                      |
|-----|-----------------------------|--------------------------------------|
| (1) | $p \vee q$                  | (Given)                              |
| (2) | $\neg q \vee r$             | (Given)                              |
| (3) | $p \vee r$                  | (Resolution of (1) and (2))          |
| (4) | $\neg p \vee r$             | (Given)                              |
| (5) | $r \vee r$                  | (Resolution of (3) and (4))          |
| (6) | $r$                         | (Idempotent law on (5))              |
| (7) | $\neg s \rightarrow \neg r$ | (Given)                              |
| (8) | $\neg \neg s$               | (Modus Tollens of (6) and (7))       |
| (9) | $\therefore s$              | (Double negative equivalence of (8)) |

# Playposit

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What rules of inference are used in the following argument?

(1)		$p \wedge q$	(Given)
(2)		$p$	(??)
(3)		$p \rightarrow r$	(Given)
(4)		$r \rightarrow s$	(Given)
(5)		$p \rightarrow s$	(??)
<hr/>			
(6)	$\therefore$	$s$	(??)

- A. Modus Ponens
- B. Modus Tollens
- C. Hypothetical Syllogism
- D. Disjunctive Syllogism
- E. Addition
- F. Simplification
- G. Conjunction
- H. Resolution

# Remember this argument?

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**Socrates is a man**  
**All men are mortal**

---

**∴ Socrates is mortal**

How do we prove this with rules of inference?



# Additional Rules of Inference for Predicates

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## 1. Universal Instantiation

- If we know something is true for the whole population (or domain  $D$ ), we can conclude that it is true for a **specific** member of the group

$$\frac{\forall x P(x), x \in D}{\therefore P(d) \text{ if } d \in D}$$

# Example

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- Consider predicates:
  - $M(x)$  :  $x$  is a man,  $x \in \text{people}$
  - $S(x)$  :  $x$  is a mortal,  $x \in \text{people}$
- And the hypotheses:
  - All men are mortal ( $\forall x M(x) \rightarrow S(x)$ ,  $x \in \text{people}$ )
  - Socrates is a man ( $M(\text{Socrates})$ )

# Example

## Propositions:

$M(x)$  :  $x$  is a man,  $x \in \text{people}$

$S(x)$  :  $x$  is a mortal,  $x \in \text{people}$

## Givens and Conclusion:

$\forall x M(x) \rightarrow S(x), x \in \text{people}$

$M(\text{Socrates})$

$\therefore ?$

- |     |  |                               |
|-----|--|-------------------------------|
| (1) | $\forall x M(x) \rightarrow S(x), x \in \text{people}$ | (Given)                       |
| (2) | $M(\text{Socrates}) \rightarrow S(\text{Socrates})$    | (Universal Instantiation)     |
| (3) | $M(\text{Socrates})$                                   | (Given)                       |
| (4) | $\therefore S(\text{Socrates})$                        | (Modus Ponens of (2) and (3)) |

# Additional Rules of Inference for Predicates

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## 2. Universal Generalization

- If we know something is true for the an arbitrary element in the  $D$  and we've made no assumptions about that element, we can conclude that it is true for all elements in the  $D$

$$\frac{P(d) \text{ for } \mathbf{arbitrary} \ d \in D}{\therefore \forall x P(x), x \in D}$$

**Caveat: Must ensure  $d$  is arbitrary!**

# Example

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- Prove  $\forall x Q(x)$  from the hypotheses:
  1.  $\forall x (P(x) \rightarrow Q(x)), x \in D$  (Given)
  2.  $\forall x P(x), x \in D$  (Given)
  3.  $P(c) \rightarrow Q(c)$ , arbitrary  $c \in D$  (Universal Instantiation)
  4.  $P(c)$  (Universal Instantiation)
  5.  $Q(c)$  (Modus Ponens of (3) and (4))

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  6.  $\therefore \forall x Q(x)$  (Universal Generalization)

# Additional Rules of Inference for Predicates

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## 3. Existential Instantiation

- We know some element, say  $d$ , such that  $P(d)$  is true.

$$\frac{\exists x P(x), x \in D}{\therefore P(d) \text{ for some } d \in D}$$

- Here,  $d$  is a new name we give that specific element
  - $d$  is **not** arbitrary, otherwise we would make false claims
  - e.g. “There exists an animal that flies, thus dogs flies.”

# Example

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- Consider the hypotheses  $\exists x P(x)$  and  $\forall x \neg P(x), x \in D$ . Show that we can derive a contradiction (i.e. **false**) from these hypotheses

- $\exists x P(x), x \in D$  (Given)
- $\forall x \neg P(x), x \in D$  (Given)
- $P(d), d$  is a specific (Existential Instantiation)  
element of  $D$
- $\neg P(d)$  (Universal Instantiation)
- $P(d) \wedge \neg P(d)$  (Conjunction of (3) and (4))

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- $\therefore$  **False** (Negation Law)

# Additional Rules of Inference for Predicates

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## 4. Existential Generalization

- If we know that  $P(d)$  is true for a specific element of  $D$ , then we know that there exists at least one element of  $D$  where  $P(x)$  is true.

$$\frac{P(d) \text{ for some } d \in D}{\therefore \exists x P(x), x \in D}$$



# Example

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- $A(x)$  :  $x$  attends UA,  $x \in \text{People}$
- $S(x)$  :  $x$  is smart,  $x \in \text{People}$
- Given  $A(\text{George})$  and  $S(\text{George})$ , prove that  $\exists x (A(x) \wedge S(x))$ 
  1.  $A(\text{George})$  (Given)
  2.  $S(\text{George})$  (Given)
  3.  $A(\text{George}) \wedge S(\text{George})$  (Conjunction of (3) and (4))

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- 4.  $\therefore \exists x (A(x) \wedge S(x)),$  (Existential Generalization)  
 $x \in \text{People}$

# Summary of Rules of Inferences for Predicates

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Name	Rule of Inference
Universal Instantiation	$\frac{\forall x P(x), x \in D}{\therefore P(d) \text{ if } d \in D}$
Universal Generalization	$\frac{P(d) \text{ for arbitrary } d \in D}{\therefore \forall x P(x), x \in D}$
Existential Instantiation	$\frac{\exists x P(x), x \in D}{\therefore P(d) \text{ for some } d \in D}$
Existential Generalization	$\frac{P(d) \text{ for some } d \in D}{\therefore \exists x P(x), x \in D}$

# Example

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- Prove that the following premises imply  $\exists x (P(x) \wedge \neg B(x)), x \in D$ 
  1.  $\exists x (C(x) \wedge \neg B(x)), x \in D$  (Given)
  2.  $\forall x (C(x) \rightarrow P(x)), x \in D$  (Given)
  3.  $C(d) \wedge \neg B(d), d \in D$  (Exist. Instantiation of (1))
  4.  $C(d)$  (Simplification of (3))
  5.  $C(d) \rightarrow P(d)$  (Univ. Instantiation of (2))
  6.  $P(d)$  (Modus Ponens of (4) and (5))
  7.  $\neg B(d)$  (Simplification of (3))
  8.  $P(d) \wedge \neg B(d)$  (Conjunction of (6) and (7))

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  9.  $\exists x (P(x) \wedge \neg B(x))$  (Exist. Generalization of (8))

# Playposit Question

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Which rule of inference is used in the following argument (assume that the domain is CS majors):

“All CS majors must take Discrete Structures. Josie is a CS major. Therefore, she must take discrete structures.”

- A. Universal Instantiation
- B. Universal Generalization
- C. Existential Instantiation
- D. Existential Generalization

# Specious Reasoning: The Bear Patrol

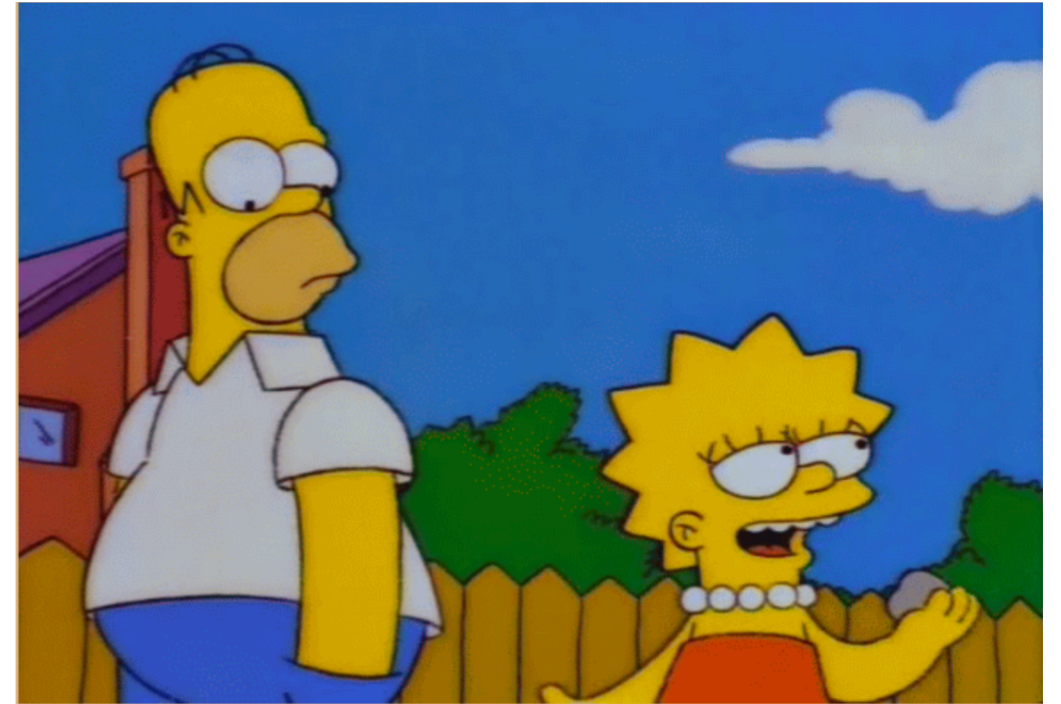
Homer: Ah, not a bear in sight. The Bear Patrol must be working like a

Lisa: That's **specious reasoning**, Dad. [...] By your logic, I could claim that this rock keeps tigers away!

Homer: Uh-huh, and how does it work?

Lisa: It doesn't work. [...] It's just a stupid rock [...] But I don't see any tigers around here, do you?

Homer: Lisa, I'd like to buy your rock.



From: **The Simpsons**, "Much Apu About Nothing"  
(Season 7, Episode 151)

# Specious Reasoning: The Bear Patrol

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## Definition: Specious Reasoning

An unsupported or improperly constructed argument.  
(That is, an unsound or invalid argument)

- Where is the error in Homer's logic?

$b$  : There are bears in Springfield

$w$  : The Bear Patrol is working

First Issue: Which of these is Homer's argument?

$$\frac{(1) \quad \neg b \quad (\text{Given})}{(2) \quad \therefore w \quad (???)}$$
$$\frac{(1) \quad w \quad (\text{Given})}{(2) \quad \therefore \neg b \quad (???)}$$

**The first seems most reasonable in this context**

# Specious Reasoning: The Bear Patrol

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- Where is the error in Homer's logic?

Next, what is the missing piece of homers argument?

(1)  $\neg b$  (Given)

(2)  $\neg b \rightarrow w$



Whoops!  
Unsupported implication

Valid  
Argument

---


(3)  $\therefore w$  (Modus Ponens from (1) and (2))

# Specious Reasoning: The Bear Patrol

---

- Where is the error in Homer's logic?

Next, what is the missing piece of homers argument?

(1)	$\neg b$	(Given)		Valid	
(2)	$\neg b \rightarrow w$			Whoops! (but unsound)	
(3)	$\therefore w$	(Modus Ponens from (1) and (2))			Argument

Ok, then how about...

(1)	$\neg b$	(Given)		Invalid	
(2)	$w \rightarrow \neg b$			Argument	
(3)	$\therefore w$	(??? from (1) and (2))			

(the second form of Homer's argument fails similarly)



# Fallacies

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**Definition:** *Fallacy* (a form of specious reasoning)

A fallacy is an argument constructed with an improper inference.

- Three classic types:
  1. Affirming the conclusion (or consequent)

Ex: If Juan is in Dallas, then he is in Texas. He is in Texas.  
Therefore, he is in Dallas.

$$\begin{array}{r} p \rightarrow q \quad (\text{Given}) \\ q \quad (\text{Given}) \\ \hline \therefore p \quad (???) \end{array}$$

# Fallacies

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## 2. Denying the Hypothesis (or Antecedent)

Ex: If Ingrid is in Dallas, then she is in Texas. She is not in Dallas. Therefore, she is not in Texas.

$$\begin{array}{r} p \rightarrow q \quad (\text{Given}) \\ \neg p \quad (\text{Given}) \\ \hline \therefore \neg q \quad (???) \end{array}$$

## 3. Begging the Question (or Circular Reasoning)

Ex: (1) I am not lying, so I must be telling the truth  
(2) The law is the law

$$\begin{array}{r} p \quad (\text{Given}) \\ \hline \therefore p \quad (???) \end{array}$$

# Playposit Question

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Which type of fallacy is the following?

“All men are mortal. Susan is not a man.  
Therefore, Susan is not mortal”

- A. Affirming the conclusion
- B. Denying the hypothesis
- C. Begging the question

# Fallacies for Fun

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## 1. Fallacy of Interrogation

Two classic examples:

- Have you stopped beating your spouse?
- Did you give your accomplice the stolen money?

## 2. No 'True Scotsman' Fallacy

“No American opposes tax cuts”

“Wendell is an American; he opposes them”

“No true American opposes tax cuts.